

# Computer-Generated Symmetric Chain Decompositions for $L(4,n)$ and $L(3,n)$

Xiangdong Wen

## 1. Introduction

Recall that the famous *Young's partition lattice*  $L(m, n)$  consists of the set of integer-vectors

$$(a_1, a_2, \dots, a_m), \quad 0 \leq a_1 \leq a_2 \leq \dots \leq a_m \leq n,$$

with the order relation

$$\vec{a} \leq \vec{b} \quad \text{if} \quad a_i \leq b_i \quad \text{for} \quad i = 1, 2, \dots, m.$$

The **rank**  $r$  is defined by

$$r(\vec{a}) = \sum_{i=1}^m a_i.$$

And recall that a chain  $\vec{v}_1 \leq \vec{v}_2 \leq \dots \leq \vec{v}_k$  in  $L(m, n)$  is called **saturated** if it skips no ranks and **symmetric** if  $r(\vec{v}_1) + r(\vec{v}_k) = mn$ .

A **Symmetric chain decomposition**(SCD) of a poset is a way of expressing it as a disjoint union of saturated symmetric chains.

One of major problems in order theory is the explicit construction of SCDs for Young's Lattice for *all*  $m$  and  $n$ . In 1989, Kathy O'Hara ([2], see also [5]) has astounded the combinatorial world by constructing SCDs for the 'trivial extension' of  $L(m, n)$ , in which all partitions of one rank are related to the next; but the problem is remaining wide open for Young's lattice itself.

SCDs for  $L(4,n)$  and  $L(3,n)$  have been constructed by West[4] and Lindström[1]. In this paper we explicitly provide **complete** SCDs for  $L(4,n)$  and  $L(3,n)$ , which were found by the assistance of our computer. And far more interestingly, it is proved *completely automatically* without using any human help (except for writing the general Maple program). While our construction for  $L(4, n)$  is not equivalent to West's construction, it is nevertheless of the similar format. On the other hand, our construction for  $L(3, n)$  is more elegant than Lindström's, since it is not split into even and odd cases.

We hope the present approach would ultimately lead to computer-generated or at least computer-assisted constructions of SCDs for  $L(m, n)$ , or at least for  $L(5, n)$ . Meanwhile we are unable to do the case of  $L(5, n)$ . We also hope the present *methodology* will be useful for future attacks on this challenging and tantalizing problem.

## 2. New SCDs for $L(3,n)$ and $L(4,n)$

**Theorem 1** *The following tables give symmetric chain decompositions for  $L(3, n)$  and  $L(4, n)$ . In tables below,  $i, j$  and  $k$  are generic non-negative integers, and vertical dots represent that the only component that is not the same gets decremented by 1. For example,  $(n - i - 3j, n - i - 2j, n - j) \dots (i + j, n - i - 2j, n - j)$  is a shorthand for the chain:*

$$(n - i - 3j - a, n - i - 2j, n - j), a = 0 \dots n - 2i - 4j.$$

$C_{ij}$ $2i + 4j \leq n$			$D_{ij}$ $2i + 4j + 3 \leq n$		
$(n-i-3j,$	$n-i-2j,$	$n-j)$	$(n-i-3j-1,$	$n-i-2j-1,$	$n-j-1)$
$\vdots$			$\vdots$		
$(i+j,$	$n-i-2j,$	$n-j)$	$(n-i-3j-1,$	$n-i-2j-1,$	$n-i-j-1)$
$\vdots$			$\vdots$		
$(i+j,$	$i+2j,$	$n-j)$	$(j,$	$n-i-2j-1,$	$n-i-j-1)$
$\vdots$			$\vdots$		
$(i+j,$	$i+2j,$	$i+3j)$	$(j,$	$i+2j+1,$	$n-i-j-1)$
$\vdots$			$\vdots$		
$(j,$	$i+2j,$	$i+3j)$	$(j,$	$i+2j+1,$	$i+3j+2)$

*A complete SCD for  $L(3,n)$*

$C_{ijk}$ $2i + 2j + 3k \leq n$				$D_{ijk}$ $2i + 2j + 3k + 3 \leq n$			
$(n-2k-2j-i,$	$n-2k-j-i,$	$n-k,j,$	$n-k)$	$(n-2k-2j-i-1,$	$n-2k-j-i-1,$	$n-k-j-1,$	$n-k)$
$\vdots$				$\vdots$			
$(k+i,$	$n-2k-j-i,$	$n-k,j,$	$n-k)$	$(n-2k-2j-i-1,$	$n-2k-j-i-1,$	$n-k-i-j-1,$	$n-k)$
$\vdots$				$\vdots$			
$(k+i,$	$k+j+i,$	$n-k,j,$	$n-k)$	$(n-2k-2j-i-1,$	$n-2k-j-i-1,$	$n-k-i-j-1,$	$n-k-i-1)$
$\vdots$				$\vdots$			
$(k+i,$	$k+j+i,$	$2k+j+i,$	$n-k)$	$(k,$	$n-2k-j-i-1,$	$n-k-i-j-1,$	$n-k-i-1)$
$\vdots$				$\vdots$			
$(k+i,$	$k+j+i,$	$2k+j+i,$	$2k+2j+i)$	$(k,$	$k+j,$	$n-k-i-j-1,$	$n-k-i-1)$
$\vdots$				$\vdots$			
$(k,$	$k+j+i,$	$2k+j+i,$	$2k+2j+i)$	$(k,$	$k+j,$	$2k+j+i+1,$	$n-k-i-1)$
$\vdots$				$\vdots$			
$(k,$	$k+j,$	$2k+j+i,$	$2k+2j+i)$	$(k,$	$k+j,$	$2k+j+i+1,$	$2k+2j+i+2)$

*A complete SCD for  $L(4,n)$*

**Proof:** The chains are clearly saturated and symmetric. Thus we only need to prove that each vector in  $L(m, n)$  ( $m = 3, 4$ ) appears only once in the tables. We introduce the commuting indeterminate  $x_1, x_2, \dots, x_m$  and  $t$ , and define the **weight** for a vector  $\vec{a} = (a_1, a_2, \dots, a_m)$  in  $L(m, n)$  as the following:

$$w(\vec{a}) = (x_0 t)^{n-a_m} (x_1 t)^{a_m - a_{m-1}} \dots (x_{m-1} t)^{a_2 - a_1} (x_m t)^{a_1}$$

For a fixed  $m$ , it is easy to see that the total weight,

$$\sum_{n=0}^{\infty} w(\vec{a}), \quad \text{where } \vec{a} \in L(m, n),$$

is a generating function

$$G(t; x_1, x_2, \dots, x_m) = \frac{1}{(1 - x_0 t)(1 - x_1 t) \dots (1 - x_m t)}.$$

Each term in the expanded power series of  $G(t; x_1, x_2, \dots, x_m)$  has coefficient 1 and corresponds to a unique vector in  $L(m, n)$ . On the other hand, for each vector in  $L(m, n)$ , there is a unique corresponding term in the power series of  $G(t; x_1, x_2, \dots, x_m)$ . Therefore, to prove that each vector in  $L(m, n)$  ( $m = 3, 4$ ) appears only once in the SCDs is the same as to prove that the total weights of the vectors in the given chains are  $G(t; x_1, x_2, x_3)$  and  $G(t; x_1, x_2, x_3, x_4)$  respectively. The part of summing over all the weights of the vectors is done by computer. ■

This method of proof can be applied to any conjectured SCDs. The difficult part is to find such decompositions, here we need

human-computer interactions by using a modified greedy algorithm. Once it is found, the verification part is purely automatic by using the Maple program `Lmn`. `Lmn` can also be used to give completely automatic proofs of the validity of Lindström's and West's constructions.

For general  $m, n$ , an explicit construction of SCDs of  $L(m, n)$  is still an open problem.

### 3. Maple package

The summation of all the weights of the vectors in the given chains is automatically done by computer. The Maple package is available at <http://www.math.temple.edu/~wen/lattice/>. After downloading the file to the local disk, type

```
>read("Lmn");
```

in the maple workspace. There is detailed on-line help on how to use the procedures in the package `Lmn`. Procedures to compute the total weights of the vectors in SCDs given by Lindström[1] and by West[4] are also included in the package `Lmn`.

### 4. Acknowledgment

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**Department of Mathematics, Temple University, Philadelphia, PA 19122. [wen@math.temple.edu](mailto:wen@math.temple.edu)**